Tensors

Shmuel Friedland Univ. Illinois at Chicago

University of Kansas, August 20, 2010

Foreword

In the past ten years, tensors again became a hot topic of research in pure and applied mathematics. In applied mathematics it is driven by data which has a few parameters. In pure math. it is quantum information theory, and multilinear algebra. There are many interesting numerical and theoretical problems that need to be resolved. Tensors are related to matrices one one hand and on the other hand are related to polynomial maps.

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To paraphrase Max Noether:

Matrices were created by God and tensors by Devil.

Ranks of 3-tensors

Basic facts.

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- Complexity.

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Characterization of tensor in $\mathbb{C}^{4\times4\times4}$ of border rank 4

scalar
$$a \in \mathbb{F}$$
, vector $\mathbf{x} = (x_1, \dots, x_n)^{\top} \in \mathbb{F}^n$, matrix $A = [a_{ij}] \in \mathbb{F}^{m \times n}$, 3-tensor $\mathcal{T} = [t_{i,j,k}] \in \mathbb{F}^{m \times n \times l}$, p-tensor $\mathcal{T} = [t_{i_1,\dots,i_p}] \in \mathbb{F}^{n_1 \times \dots \times n_p}$

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Abstractly $\mathbb{U}:=\mathbb{U}_1\otimes\mathbb{U}_2\otimes\mathbb{U}_3$ dim $\mathbb{U}_i=m_i, i=1,2,3,$ dim $\mathbb{U}=m_1m_2m_3$

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Tensor calculus 1890 G. Ricci-Curbastro: absolute differential calculus,

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Rank one tensor $t_{i,j,k} = x_i y_j z_k$, $(i,j,k) = (1,1,1), \dots, (m_1, m_2, m_3)$ or decomposable tensor $\mathbf{x} \otimes \mathbf{y} \otimes \mathbf{z}$

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basis of \mathbb{U}_{j} : $[\mathbf{u}_{1,j},\dots,\mathbf{u}_{m_{j},j}] \ j=1,2,3$ basis of \mathbb{U} : $\mathbf{u}_{i_{1},1}\otimes\mathbf{u}_{i_{2},2}\otimes\mathbf{u}_{i_{3},3}, i_{j}=1,\dots,m_{j}, j=1,2,3,$

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Unfolding tensor: in direction 1:

$$\mathcal{T} = [t_{i,j,k}]$$
 view as a matrix $A_1 = [t_{i,(j,k)}] \in \mathbb{F}^{m_1 \times (m_2 \cdot m_3)}$

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 (CANDEC, PARFAC)



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Basic facts

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COR rank $T \leq \min(mn, ml, nl)$

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PRF: 3-sat with n variables m clauses satisfiable iff rank $\mathcal{T}=4n+2m, \mathcal{T}\in\mathbb{F}^{(2n+3m)\times(3n)\times(3n+m)})$ otherwise rank is larger

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In all the examples we know $\operatorname{mtrank}(m, n, l) \leq \operatorname{grank}(m, n, l) + 1$

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Reason: For I = (m-1)(n-1) + 1 a generic subspace of matrices of dimension I in $\mathbb{C}^{m \times n}$ intersect the variety of rank one matrices in $\mathbb{C}^{m \times n}$ at least at I lines which contain I linearly independent matrices

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$$\operatorname{Conjecture} \operatorname{grank}_{\mathbb{C}}(m, n, l) = \lceil \frac{mnl}{(m + n + l - 2)} \rceil$$

$$\operatorname{for} 2 \leq m \leq n \leq l < (m - 1)(n - 1) \text{ and } (3, n, l) \neq (3, 2p + 1, 2p + 1)$$

THM: $\operatorname{grank}_{\mathbb{C}}(m, n, I) = \min(I, mn)$ for $(m-1)(n-1) + 1 \leq I$.

Reason: For I = (m-1)(n-1) + 1 a generic subspace of matrices of dimension I in $\mathbb{C}^{m \times n}$ intersect the variety of rank one matrices in $\mathbb{C}^{m \times n}$ at least at I lines which contain I linearly independent matrices

COR: grank_C
$$(2, n, I) = \min(I, 2n)$$
 for $2 \le n \le I$

$$f_r: (\mathbb{C}^m \times \mathbb{C}^n \times \mathbb{C}^l)^r \to \mathbb{C}^{m \times n \times l}, \mathbf{x} \otimes \mathbf{y} \otimes \mathbf{z} = (a\mathbf{x}) \otimes (b\mathbf{y}) \otimes ((ab)^{-1}\mathbf{z})$$

$$\operatorname{grank}_{\mathbb{C}}(m, n, l)(m + n + l - 2) \geq mnl \Rightarrow \operatorname{grank}_{\mathbb{C}}(m, n, l) \geq \lceil \frac{mnl}{(m + n + l - 2)} \rceil$$

$$\operatorname{Conjecture} \operatorname{grank}_{\mathbb{C}}(m, n, l) = \lceil \frac{mnl}{(m + n + l - 2)} \rceil$$

$$\operatorname{for} 2 \leq m \leq n \leq l < (m - 1)(n - 1) \text{ and } (3, n, l) \neq (3, 2p + 1, 2p + 1)$$

$$\operatorname{Fact:} \operatorname{grank}_{\mathbb{C}}(3, 2p + 1, 2p + 1) = \lceil \frac{3(2p + 1)^2}{4p + 3} \rceil + 1$$

bilinear map: $\phi : \mathbf{U} \times \mathbf{V} \to \mathbf{W}$

 $[\mathbf{u}_1,\ldots,\mathbf{u}_m],[\mathbf{v}_1,\ldots,\mathbf{v}_n],[\mathbf{w}_1,\ldots,\mathbf{w}_l]$ bases in $\mathbf{U},\mathbf{V},\mathbf{W}$

$$[\mathbf{u}_1,\ldots,\mathbf{u}_m],[\mathbf{v}_1,\ldots,\mathbf{v}_n],[\mathbf{w}_1,\ldots,\mathbf{w}_l]$$
 bases in $\mathbf{U},\mathbf{V},\mathbf{W}$

$$\phi(\mathbf{u}_i, \mathbf{v}_j) = \sum_{k=1} t_{i,j,k} \mathbf{w}_k, \, \mathcal{T} := [t_{i,j,k}] \in \mathbb{F}^{m \times n \times l}$$

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$$\mathcal{T} = \sum_{a=1}^{r} \mathbf{x}_{a} \otimes \mathbf{y}_{a} \otimes \mathbf{z}_{a}, r = \mathrm{rank} \ \mathcal{T}$$

$$\phi(\mathbf{c}, \mathbf{d}) = \sum_{a=1}^{r} (\mathbf{c}^{\top} \mathbf{x}) (\mathbf{d}^{\top} \mathbf{y}) \mathbf{z}_{a}, \ \mathbf{c} = \sum_{i=1}^{m} c_{i} \mathbf{u}_{i}, \mathbf{d} = \sum_{j=1}^{n} d_{j} \mathbf{v}_{j}$$

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Complexity: r -products

bilinear map: $\phi : \mathbf{U} \times \mathbf{V} \to \mathbf{W}$

$$[\mathbf{u}_1,\ldots,\mathbf{u}_m],[\mathbf{v}_1,\ldots,\mathbf{v}_n],[\mathbf{w}_1,\ldots,\mathbf{w}_l]$$
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Complexity: r -products

Matrix product $\tau : \mathbb{F}^{M \times N} \times \mathbb{F}^{N \times L} \to \mathbb{F}^{M \times L}$, $(A, B) \mapsto AB$



$$[\mathbf{u}_1,\ldots,\mathbf{u}_m],[\mathbf{v}_1,\ldots,\mathbf{v}_n],[\mathbf{w}_1,\ldots,\mathbf{w}_l]$$
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Matrix product
$$\tau : \mathbb{F}^{M \times N} \times \mathbb{F}^{N \times L} \to \mathbb{F}^{M \times L}$$
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$$M=N=L=2$$
, $\operatorname{grank}_{\mathbb{C}}(4,4,4)=\lceil \frac{4\times 4\times 4}{4+4+4-2} \rceil = \lceil 6.4 \rceil = 7$



bilinear map: $\phi : \mathbf{U} \times \mathbf{V} \to \mathbf{W}$

$$[\mathbf{u}_1,\ldots,\mathbf{u}_m],[\mathbf{v}_1,\ldots,\mathbf{v}_n],[\mathbf{w}_1,\ldots,\mathbf{w}_l]$$
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Complexity: r -products

Matrix product
$$\tau : \mathbb{F}^{M \times N} \times \mathbb{F}^{N \times L} \to \mathbb{F}^{M \times L}$$
, $(A, B) \mapsto AB$

$$M=N=L=2$$
, grank_C $(4,4,4)=\lceil \frac{4\times4\times4}{4+4+4-2} \rceil = \lceil 6.4 \rceil = 7$

Product of two 2×2 matrices is done by 7 multiplications



$$\text{grank}(3,2p,2p) = \lceil \tfrac{12p^2}{4p+1} \rceil \text{ and } \text{grank}(3,2p-1,2p-1) = \lceil \tfrac{3(2p-1)^2}{4p-1} \rceil + 1$$

grank
$$(3, 2p, 2p) = \lceil \frac{12p^2}{4p+1} \rceil$$
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```
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```

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\begin{aligned} & \operatorname{grank}(3,2p,2p) = \lceil \frac{12p^2}{4p+1} \rceil \text{ and } \operatorname{grank}(3,2p-1,2p-1) = \lceil \frac{3(2p-1)^2}{4p-1} \rceil + 1 \\ & (n,n,n+2) \text{ if } n \neq 2 \pmod{3}, \\ & (n-1,n,n) \text{ if } n = 0 \pmod{3}, \\ & (4,m,m) \text{ if } m \geq 4, \\ & (n,n,n) \text{ if } n \geq 4 \\ & (l,2p,2q) \text{ if } l \leq 2p \leq 2q \text{ and and } \frac{2lp}{l+2p+2q-2} \text{ is integer} \end{aligned}
```

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```

Easy to compute grank_{\mathbb{C}}(m, n, l):

$$\begin{array}{l} \operatorname{grank}(3,2p,2p) = \lceil \frac{12p^2}{4p+1} \rceil \text{ and } \operatorname{grank}(3,2p-1,2p-1) = \lceil \frac{3(2p-1)^2}{4p-1} \rceil + 1 \\ (n,n,n+2) \text{ if } n \neq 2 \pmod{3}, \\ (n-1,n,n) \text{ if } n = 0 \pmod{3}, \\ (4,m,m) \text{ if } m \geq 4, \\ (n,n,n) \text{ if } n \geq 4 \\ (l,2p,2q) \text{ if } l \leq 2p \leq 2q \text{ and and } \frac{2lp}{l+2p+2q-2} \text{ is integer} \end{array}$$

Easy to compute grank_{\mathbb{C}}(m, n, l):

Pick at random
$$\mathbf{w}_r := (\mathbf{x}_1, \mathbf{y}_1, \mathbf{z}_1, \dots, \mathbf{x}_r, \mathbf{y}_r, \mathbf{z}_r) \in (\mathbb{R}^m \times \mathbb{R}^n \times \mathbb{R}^l)^r$$

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The minimal $r \geq \lceil \frac{mnl}{(m+n+l-2)} \rceil$ s.t. rank $J(f_r)(\mathbf{w}_r) = mnl$ is $\operatorname{grank}_{\mathbb{C}}(m, n, l)$ (Terracini Lemma 1915)

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$$\mathbf{w}_r := (\mathbf{x}_1, \mathbf{y}_1, \mathbf{z}_1, \dots, \mathbf{x}_r, \mathbf{y}_r, \mathbf{z}_r) \in (\mathbb{R}^m \times \mathbb{R}^n \times \mathbb{R}^l)^r$$

The minimal $r \geq \lceil \frac{mnl}{(m+n+l-2)} \rceil$ s.t. rank $J(f_r)(\mathbf{w}_r) = mnl$ is $\operatorname{grank}_{\mathbb{C}}(m, n, l)$ (Terracini Lemma 1915)

Avoid round-off error:

 $\mathbf{w}_r \in (\mathbb{Z}^m \times \mathbb{Z}^n \times \mathbb{Z}^l)^r$ find rank $J(f_r)(\mathbf{w}_r)$ exact arithmetic



```
\begin{array}{l} \operatorname{grank}(3,2p,2p) = \lceil \frac{12p^2}{4p+1} \rceil \text{ and } \operatorname{grank}(3,2p-1,2p-1) = \lceil \frac{3(2p-1)^2}{4p-1} \rceil + 1 \\ (n,n,n+2) \text{ if } n \neq 2 \pmod{3}, \\ (n-1,n,n) \text{ if } n = 0 \pmod{3}, \\ (4,m,m) \text{ if } m \geq 4, \\ (n,n,n) \text{ if } n \geq 4 \\ (l,2p,2q) \text{ if } l \leq 2p \leq 2q \text{ and and } \frac{2lp}{l+2p+2q-2} \text{ is integer} \end{array}
```

Easy to compute grank_{\mathbb{C}}(m, n, l):

Pick at random
$$\mathbf{w}_r := (\mathbf{x}_1, \mathbf{y}_1, \mathbf{z}_1, \dots, \mathbf{x}_r, \mathbf{y}_r, \mathbf{z}_r) \in (\mathbb{R}^m \times \mathbb{R}^n \times \mathbb{R}^l)^r$$

The minimal $r \geq \lceil \frac{mnl}{(m+n+l-2)} \rceil$ s.t. rank $J(f_r)(\mathbf{w}_r) = mnl$ is $\operatorname{grank}_{\mathbb{C}}(m, n, l)$ (Terracini Lemma 1915)

Avoid round-off error:

 $\mathbf{w}_r \in (\mathbb{Z}^m \times \mathbb{Z}^n \times \mathbb{Z}^l)^r$ find rank $J(f_r)(\mathbf{w}_r)$ exact arithmetic I checked the conjecture up to $m, n, l \leq 14$

For $mn \le I$ mtrank(m, n, I) = grank(m, n, I) = mn.

For $mn \le I$ mtrank(m, n, I) = grank(m, n, I) = mn.

For $mn \le l$ mtrank(m, n, l) = grank(m, n, l) = mn.

Closure
$$(\bigcup_{i=1}^{c(m,n,l)}) = \mathbb{R}^{m \times n \times l}$$

For $mn \le I$ mtrank(m, n, I) = grank(m, n, I) = mn.

Closure(
$$\bigcup_{i=1}^{c(m,n,l)}$$
) = $\mathbb{R}^{m \times n \times l}$ rank $\mathcal{T} = \operatorname{grank}(m,n,l)$ for each $\mathcal{T} \in V_1$

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Closure
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rank $\mathcal{T} = \operatorname{grank}(m,n,l)$ for each $\mathcal{T} \in V_1$
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```
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\{\rho_1, \dots, \rho_{c(m,n,l)}\} = \{\operatorname{grank}(m,n,l), \dots, \operatorname{mtrank}(m,n,l)\}
```

For $mn \le l$ mtrank(m, n, l) = grank(m, n, l) = mn.

For $2 \le m \le n \le l < mn-1$, there exist $V_1, \ldots, V_{c(m,n,l)} \subset \mathbb{R}^{m \times n \times l}$ pairwise distinct open connected semi-algebraic sets s.t.

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 $\operatorname{mtrank}(2, n, I) = \operatorname{grank}(2, n, I) = \min(I, 2n) \text{ if } 2 \le n < I \text{ - one typical rank}$ $\operatorname{mtrank}(2, n, n) = \operatorname{grank}(2, n, n) + 1 = n + 1 \text{ if } 2 \le n \text{ - two typical ranks}$

For $mn \le I$ mtrank(m, n, I) = grank(m, n, I) = mn.

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For
$$I = (m-1)(n-1) + 1 \exists m, n$$
:
 $c(m, n, l) > 1$, mtrank $(m, n, l) \ge \text{grank}(m, n, l) + 1$

For $mn \le I$ mtrank(m, n, I) = grank(m, n, I) = mn.

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Examples [3]



$$\mathbb{R}^{m imes n imes l}$$
 IPS: $\langle \mathcal{A}, \mathcal{B}
angle = \sum_{i=j=k}^{m,n,l} a_{i,j,k} b_{i,j,k}, \ \|\mathcal{T}\| = \sqrt{\langle \mathcal{T}, \mathcal{T}
angle}$

$$\begin{array}{l} \mathbb{R}^{m\times n\times l} \ \text{IPS:} \ \langle \mathcal{A},\mathcal{B}\rangle = \sum_{i=j=k}^{m,n,l} a_{i,j,k} b_{i,j,k}, \ \|\mathcal{T}\| = \sqrt{\langle \mathcal{T},\mathcal{T}\rangle} \\ \langle \mathbf{x}\otimes\mathbf{y}\otimes\mathbf{z},\mathbf{u}\otimes\mathbf{v}\otimes\mathbf{w}\rangle = (\mathbf{u}^{\top}\mathbf{x})(\mathbf{v}^{\top}\mathbf{y})(\mathbf{w}^{\top}\mathbf{z}) \end{array}$$

$$\begin{array}{l} \mathbb{R}^{m\times n\times l} \ \text{IPS:} \ \langle \mathcal{A},\mathcal{B} \rangle = \sum_{i=j=k}^{m,n,l} a_{i,j,k} b_{i,j,k}, \ \|\mathcal{T}\| = \sqrt{\langle \mathcal{T},\mathcal{T} \rangle} \\ \langle \mathbf{x} \otimes \mathbf{y} \otimes \mathbf{z}, \mathbf{u} \otimes \mathbf{v} \otimes \mathbf{w} \rangle = (\mathbf{u}^{\top} \mathbf{x}) (\mathbf{v}^{\top} \mathbf{y}) (\mathbf{w}^{\top} \mathbf{z}) \end{array}$$

X subspace of $\mathbb{R}^{m \times n \times l}$, $\mathcal{X}_1, \dots, \mathcal{X}_d$ an orthonormal basis of **X**

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X subspace of $\mathbb{R}^{m \times n \times l}$, $\mathcal{X}_1, \dots, \mathcal{X}_d$ an orthonormal basis of **X** $P_{\mathbf{X}}(\mathcal{T}) = \sum_{i=1}^{d} \langle \mathcal{T}, \mathcal{X}_i \rangle \mathcal{X}_i$, $\|P_{\mathbf{X}}(\mathcal{T})\|^2 = \sum_{i=1}^{d} \langle \mathcal{T}, \mathcal{X}_i \rangle^2$

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 λ singular value, $\mathbf{x}, \mathbf{y}, \mathbf{z}$ singular vectors



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How many distinct singular values are for a generic tensor?



ℓ_p maximal problem and Perron-Frobenius

$$\|(x_1,\ldots,x_n)^{\top}\|_{p}:=(\sum_{i=1}^n|x_i|^p)^{\frac{1}{p}}$$

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p = 3 is most natural in view of homogeneity

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Assume that $T \ge 0$. Then $\mathbf{x}, \mathbf{y}, \mathbf{z} \ge 0$

For which values of *p* we have an analog of Perron-Frobenius theorem?

Yes, for $p \ge 3$, No, for p < 3, Friedland-Gauber-Han [1]



 (R_1, R_2, R_3) -rank approximation of 3-tensors

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Fundamental problem in applications:

Approximate well and fast $T \in \mathbb{R}^{m_1 \times m_2 \times m_3}$ by rank (R_1, R_2, R_3) 3-tensor.

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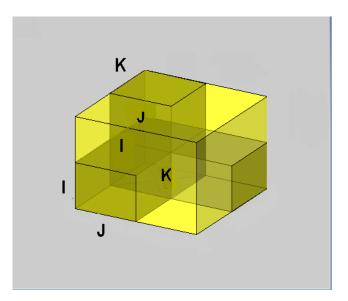
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Use Newton method on Grassmannians - Eldén-Savas 2009 [1]



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 $\mathsf{min}_{\mathcal{U} \in \mathbb{C}^{p \times q \times r}} \, \|\mathcal{A} - \mathcal{U} \times F \times E \times G\|_F \text{ achieved for } \mathcal{U} = \mathcal{A} \times E^\dagger \times F^\dagger \times G^\dagger$

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$$\min_{U \in \mathbb{C}^{p \times q}} \|A - CUR\|_F$$
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Faster choice: $U = A[I, J]^{\dagger}$ (corresponds to best *CUR* approximation on the entries read)

For given
$$A \in \mathbb{R}^{m \times n \times I}$$
, $F \in \mathbb{R}^{m \times p}$, $E \in \mathbb{R}^{n \times q}$, $G \in \mathbb{R}^{I \times r}$, where $\langle p \rangle \subset \langle n \rangle \times \langle I \rangle$, $\langle q \rangle \subset \langle m \rangle \times \langle I \rangle$, $\langle r \rangle \subset \langle m \rangle \times \langle I \rangle$

$$\min_{\mathcal{U} \in \mathbb{C}^{p \times q \times r}} \|\mathcal{A} - \mathcal{U} \times F \times E \times G\|_F$$
 achieved for $\mathcal{U} = \mathcal{A} \times E^\dagger \times F^\dagger \times G^\dagger$

CUR approximation of A obtained by choosing E, F, G submatrices of unfolded A in the mode 1, 2, 3.

Face recognition

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Video tracking

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Factor analysis

$$0 \le \mathcal{T} = [t_{i,j,k}] \in \mathbb{R}^{m \times n \times l}$$
 has given row, column and depth sums: $\mathbf{r} = (r_1, \dots, r_m)^\top, \mathbf{c} = (c_1, \dots, c_n)^\top, \mathbf{d} = (d_1, \dots, d_l)^\top > \mathbf{0}$:

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Solution: Convert to the minimal problem:

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 f_T is convex

 $f_{\mathcal{T}}$ is strictly convex implies \mathcal{T} is not decomposable: $\mathcal{T} \neq \mathcal{T}_1 \oplus \mathcal{T}_2$.

if f_T is strictly convex and is ∞ on ∂S , f_T achieves its unique minimum

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Equivalent to: the inequalities $x_i + y_j + z_k \le 0$ if $t_{i,j,k} > 0$ and equalities $\mathbf{r}^\top \mathbf{x} = \mathbf{c}^\top \mathbf{y} = \mathbf{d}^\top \mathbf{z} = 0$ imply $\mathbf{x} = \mathbf{0}_m, \mathbf{y} = \mathbf{0}_n, \mathbf{z} = \mathbf{0}_l$.

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Fact: For $\mathbf{r} = \mathbf{1}_m$, $\mathbf{c} = \mathbf{1}_n$, $\mathbf{d} = \mathbf{1}_l$ Sinkhorn scaling algorithm works.

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Major problem in algebraic statistics: phylogenic trees and their invariants [1]

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Friedland [5] one needs a equations of degree 16

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