

Balancing Search Trees

1 Tree Balance and Rotation

- binary search trees
- right rotation of a tree around a node
- code for right rotation

2 AVL Trees

- self-balancing search trees
- four kinds of critically unbalanced trees

3 code for rotation

- from left-right to left-left tree

MCS 360 Lecture 33
Introduction to Data Structures
Jan Vershelde, 13 April 2020

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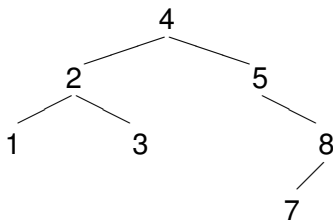
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Binary Search Trees

Consider 4, 5, 2, 3, 8, 1, 7 (recall lecture 24).

Insert the numbers in a tree:



Rules to insert x at node N :

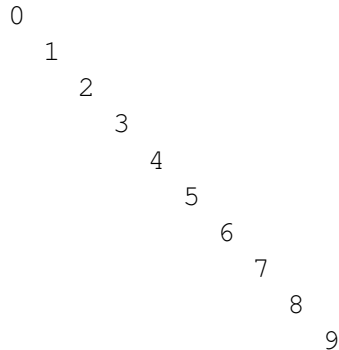
- if N is empty, then put x in N
- if $x < N$, insert x to the left of N
- if $x \geq N$, insert x to the right of N

Recursive printing: left, node, right sorts the sequence.

an unbalanced tree

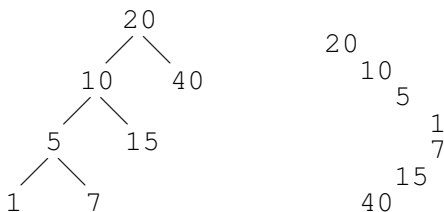
Inserting 0, 1, 2, ..., 9.

depth of tree : 9



shaping binary search trees

To make a binary search tree with given shape:



Insert numbers in a particular order: 20, 40, 10, 5, 15, 1, 7.

$$\begin{aligned} \text{depth}(T) &= 0, \text{ if } T \text{ is empty,} \\ &= 1 + \max(\text{depth}(\text{left}(T)), \text{depth}(\text{right}(T))), \text{ otherwise.} \end{aligned}$$

The tree is unbalanced because the depth of the left tree is two, while the depth of the right tree is zero.

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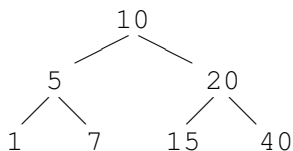
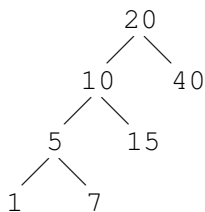
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Right Rotation

To balance the binary search tree, we do a right rotate around the root:

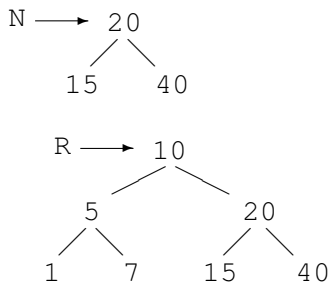
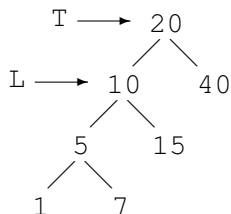


Observe the effects of a right rotation:

- left tree has become the new root;
- old root is now at the right of new root;
- left tree of old root is now the right tree of the left tree of old root.

Right Rotation in 3 Steps

Tree with root node T:



- 1 Label left of T with L.
- 2 New tree N has right of T as right and as left the right of L.
- 3 Result R has L as root, the tree N as right, and the left of L as left.

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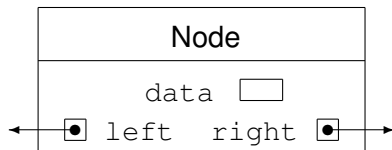
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a node struct



```
struct Node
{
    int data;    // numbers stored at node in tree
    Node *left; // pointer to left branch of tree
    Node *right; // pointer to right branch of tree

    Node(const int& item, Node* left_ptr = NULL,
          Node* right_ptr = NULL) :
        data(item),
        left(left_ptr), right(right_ptr) {}
}
```

a class Tree

```
#include "mcs360_integer_tree_node.h"

namespace mcs360_integer_tree
{
    class Tree
    {
    private:
        Node *root; // data member

    public:
        Tree(const int& item,
            const Tree& left = Tree(),
            const Tree& right = Tree() ) :
            root(new Node(item, left.root, right.root)) {}
        Tree get_left() const;
        Tree get_right() const;
        void insert(int item);
    };
}
```

function rotate_right

Prototype of function in client of class Tree:

```
Tree rotate_right ( Tree t );  
  
// Returns the tree rotated to the right  
// around its root.  
  
// Precondition: left of t is not null.
```

definition of rotate_right

```
Tree rotate_right ( Tree t )
{
    Tree left = t.get_left();

    Tree new_t = Tree(t.get_data(),
        left.get_right(),t.get_right());

    Tree R = Tree(left.get_data(),
        left.get_left(),new_t);

    return R;
}
```

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AVL Trees

Define the balance of a tree as

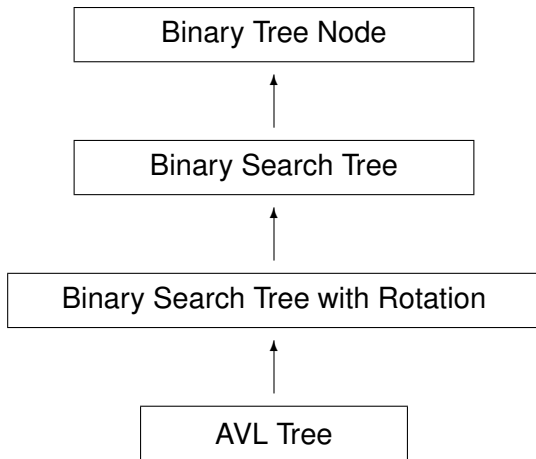
$$\text{balance} = \text{depth}(\text{right tree}) - \text{depth}(\text{left tree}).$$

G.M. Adel'son-Vel'skiî and E.M Landis published in 1962 an algorithm to maintain the balance of a binary search tree.

If balance gets out of range $-1 \dots + 1$, the subtree is rotated to bring into balance.

Their approach is known as *AVL trees*.

a Class Hierarchy

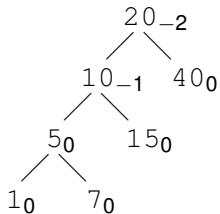


computing the balance

Recall the definition:

$$\text{balance} = \text{depth}(\text{right tree}) - \text{depth}(\text{left tree}).$$

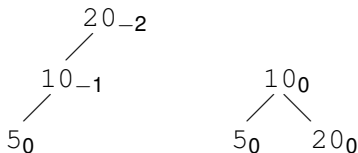
At every node we compute the balance, displayed as subscript:



balancing a left-left tree

The tree below is *left heavy* as the balance is -2 .

We also say that this is a *left-left tree*.



Executing a right rotation balances the tree.

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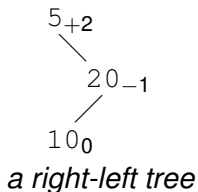
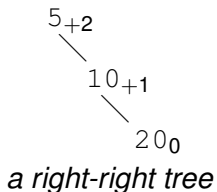
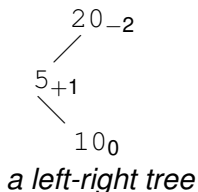
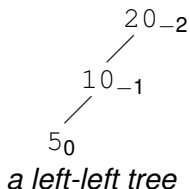
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critically unbalanced trees

A tree is *critically unbalanced* if its balance is -2 or $+2$.

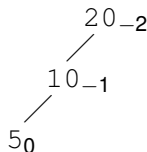
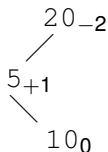


balancing trees of mixed kind

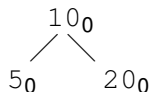
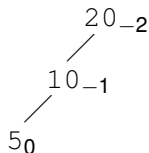
A right rotation balances a left-left tree
and a left rotation balances a right-right tree.

Balancing a left-right tree happens in two stages:

1 rotate left-right tree to left-left tree:

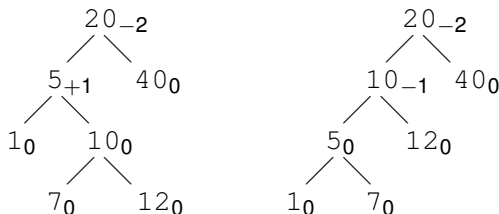


2 apply right rotation to left-left tree:



rotating a left-right tree

We rotate the left-right tree to a left-left tree:

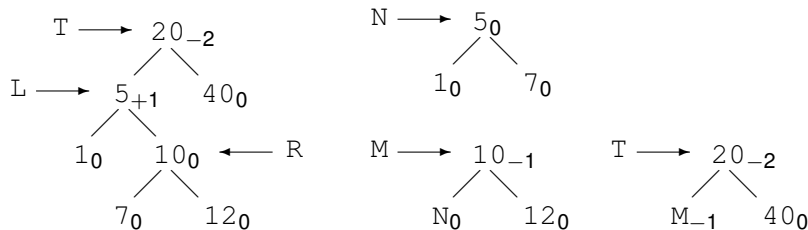


Observe the effects of the rotation:

- the data at the left node of the new tree (10) is swapped with the data of the left of the old tree (5);
- the right of the left of the new tree (12) is the right of the right of the left of the old tree;
- the right of the left of the left of the new tree (7) is the left of the right of the left of the old tree.

rotating to left-left tree in 4 steps

Tree with root node T:



- 1 Label left of T with L and right of L with R.
- 2 Tree N has as its left the left of L, as its right the left of R.
- 3 Tree M has as its left N, as its right the right of R.
- 4 Return the tree with its left M and its right the right of T.

a function to rotate a tree

```
Tree balance_by_rotation ( Tree t )
{
    if(is_left_left(t))
        return rotate_right(t);
    else if(is_right_right(t))
        return rotate_left(t);
    else if(is_left_right(t))
    {
        Tree R = rotate_to_left_left(t);
        return rotate_right(R);
    }
    else if(is_right_left(t))
    {
        Tree R = rotate_to_right_right(t);
        return rotate_left(R);
    }
}
```


the recursive calls

```
else
{
    Tree L,R;
    if(!t.is_left_null())
        L = balance_by_rotation(t.get_left());
    if(!t.is_right_null())
        R = balance_by_rotation(t.get_right());
    return Tree(t.get_data(),L,R);
}
}
```

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prototype of the function

```
Tree rotate_to_left_left ( Tree t );  
  
// Returns the tree rotated to a left-left tree.  
  
// Preconditions:  
//   (1) left of t is not null; and  
//   (2) right of left of t is not null.
```

Test: insert 20, 5, 1, 10, 7, 12 to binary search tree.

definition of the function

```
Tree rotate_to_left_left ( Tree t )
{
    Tree left = t.get_left();
    Tree right = left.get_right();

    Tree new_left = Tree(left.get_data(),
        left.get_left(), right.get_left());

    Tree new_right = Tree(right.get_data(),
        new_left, right.get_right());

    Tree R = Tree(t.get_data(),
        new_right, t.get_right());

    return R;
}
```

rebalancing search trees

After each insert (or removal):

- check the balance of the tree,
- and if critically unbalanced, rebalance.

Performance of the AVL tree:

- worst case: $1.44 \times \log_2(n)$,
- on average: $\log_2(n) + 0.25$ comparisons needed.

→ close to complete binary search tree.

Summary + Exercises

Started chapter 11 on balancing binary search trees.

Exercises:

- 1 Take a numerical example to left rotate a binary search tree with integer values. Formulate carefully each step in the left rotation. Justify the correctness of the algorithm.
- 2 Formulate the algorithm to rotate a right-left tree to a right-right tree and illustrate with an example.
- 3 The posted code provides functions to make an AVL tree. Design a class to represent an AVL tree.
- 4 Take your design of the previous exercise and define the methods of the class to represent an AVL tree.