

# Graphs

## 1 Graphs

- directed and undirected graphs
- weighted graphs
- adjacency matrices

## 2 Graph Representations

- abstract data type
- adjacency list
- adjacency matrix

## 3 Graph Implementations

- adjacency matrix
- adjacency list

MCS 360 Lecture 38  
Introduction to Data Structures  
Jan Verschelde, 27 November 2017

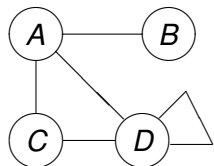
# graphs, terminology and definition

- A node in a graph is called a *vertex*.
- A connection between two vertices is an *edge*.

A *graph*  $G$  is defined by a tuple  $(V, E)$ :

- $V$  is the set of vertices, and
- $E$  is the set of edges.

An example:



$$V = \{A, B, C, D\}$$

$$E = \{(A, B), (A, C), (A, D), (C, D), (D, D)\}$$

Two vertices are *adjacent* if there is an edge between them.

A *path* is a sequence of successively adjacent vertices.

For example, a path from  $B$  to  $D$  is  $B, A, D$ .

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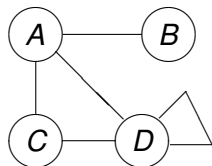
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- adjacency matrix
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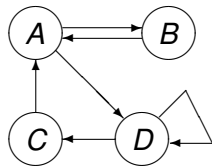
## directed and undirected graphs

In an *undirected* graph (which is the default), the edge  $(A, B)$  means: there is an edge from  $A$  to  $B$  and there is an edge from  $B$  to  $A$ .



$$V = \{A, B, C, D\}$$

$$E = \{(A, B), (A, C), (A, D), (C, D), (D, D)\}$$



$$V = \{A, B, C, D\}$$

$$E = \{(A, B), (B, A), (C, A), (A, D), (D, C), (D, D)\}$$

In a *directed* graph or *digraph*,  $(A, B)$  means there is an edge from  $A$  to  $B$  and  $(B, A)$  means there is an edge from  $B$  to  $A$ .

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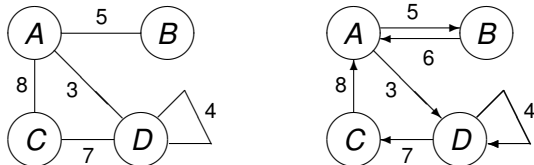
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# weighted graphs

In an *weighted* graph, a value (a weight) is associated to every edge.



Edges in a weighted graphs are triplets  $(u, v, w)$ , with  $w$  the weight of the edge which connects  $u$  to  $v$ .

A *cycle* is a path where the first vertex equals the last one. For example: A, D, C, A is a cycle.

A graph without cycles can be viewed as a tree.

The *degree* of a vertex is the number of edges that contain that vertex.

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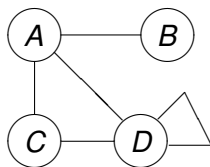
## 3 Graph Implementations

- adjacency matrix
- adjacency list

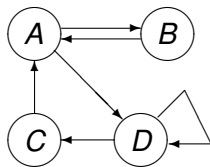
## adjacency matrices

For a graph  $G = (V, E)$ , the *adjacency matrix*  $M$

- has as many rows and columns as the number of vertices in  $V$ ;
- for all  $v_i, v_j \in V$ : if  $(v_i, v_j) \notin E$ :  $M_{i,j} = 0$ ;
- for all  $v_i, v_j \in V$ : if  $(v_i, v_j) \in E$ :  $M_{i,j} = 1$ .



	A	B	C	D
A	0	1	1	1
B	1	0	0	0
C	1	0	0	1
D	1	0	1	1

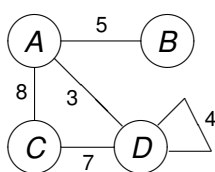


	A	B	C	D
A	0	1	0	1
B	1	0	0	0
C	1	0	0	0
D	0	0	1	1

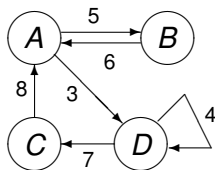
# adjacency matrices for weighted graphs

For a weighted graph  $G = (V, E)$ , the *adjacency matrix*  $M$

- has as many rows and columns as the number of vertices in  $V$ ;
- for all  $v_i, v_j \in V$ : if  $(v_i, v_j) \notin E$ :  $M_{i,j} = 0$ ;
- for all  $v_i, v_j \in V$ : if  $(v_i, v_j) \in E$ :  $M_{i,j} = w$ ,  $w$  is the weight of  $(v_i, v_j)$ .



	A	B	C	D
A	0	5	8	3
B	5	0	0	0
C	8	0	0	7
D	3	0	7	4



	A	B	C	D
A	0	5	0	3
B	6	0	0	0
C	8	0	0	0
D	0	0	7	4

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# graph representations

A graph  $G = (V, E)$  can be one of the four types:

- 1 undirected: the edge  $(u, v)$  goes from  $u$  to  $v$  and from  $v$  to  $u$ ;
- 2 directed: the edge  $(u, v)$  goes only from  $u$  to  $v$ , not from  $v$  to  $u$ ;
- 3 weighted undirected: every edge is a triplet  $(u, v, w)$ ,  
 $w$  is the weight and  $(u, v)$  goes from  $u$  to  $v$  and from  $v$  to  $u$ ;
- 4 weighted directed: every edge is a triplet  $(u, v, w)$ ,  
 $w$  is the weight and  $(u, v)$  goes only from  $u$  to  $v$ , not from  $v$  to  $u$ .

A graph  $G = (V, E)$  is a data structure:

- The set of vertices is finite:  $V = \{v_0, v_1, \dots, v_{n-1}\}$ .  
The vertices store values, the data are stored in  $V$ .
- The set of edges is finite:  $E = \{e_0, e_1, \dots, e_{m-1}\}$ .  
In a weighted graph, the edges have weights.

# operations on graphs

On a graph  $G = (V, E)$ , we want the following operations:

- Get the number  $|V|$  of vertices in  $V$ .
- Get the number  $|E|$  of edges in  $E$ .
- For an iterator  $i$  or index  $i$ , get the  $i$ -th vertex in  $V$ .
- For an iterator  $i$  or index  $i$ , get the  $i$ -th edge in  $E$ .
- Given  $u, v \in V$ , is  $(u, v) \in E$ ?
- Add a vertex to  $V$ .
- Given  $u, v \in V$ , add  $(u, v)$  to  $E$ .

For a weighted graph, an additional operation is the following:

- Given  $u, v \in V$ , get the weight of the edge from  $u$  to  $v$ .

## abstract data type

$V = \{v_0, v_1, \dots, v_{n-1}\}$  is represented by the sequence  $0, 1, \dots, n - 1$ .

```
abstract <typename T> graph;
/* A graph has vertices 0, 1, .. to store data of
   type T and edges which connect the vertices. */

abstract int number_of_vertices ( graph g );
postcondition: number_of_vertices(g) = |vertices of g|;

abstract int number_of_edges ( graph g );
postcondition: number_of_edges(g) = |edges of g|;

abstract T data_element ( graph g, size_t i );
precondition: 0 <= i < number_of_vertices(g);
postcondition: data_element(g, i) is the data at i;
```

## abstract data type continued

```
abstract (i, j) get_edge ( graph g, size_t k );  
precondition: 0 <= k < number_of_edges(g);  
postcondition: get_edge(g, k) is the k-th edge (i, j);
```

```
abstract bool is_edge ( graph g, size_t i, size_t j );  
postcondition: is_edge(g, i, j) is true if (i, j) is  
    an edge of g, otherwise is_edge(g, i, j) is false;
```

```
abstract void add_vertex ( graph g );  
postcondition: for the graph g on input, there is  
    a vertex with index == number_of_vertices(g);
```

```
abstract void add_edge ( graph g, size_t i, size_t j );  
precondition: 0 <= i, j < number_of_vertices(g);  
postcondition: is_edge(g, i, j) is true;
```

## abstract data type for weighted graph

The ADT for a weighted graph starts differently:

```
abstract <typename T> weighted graph;  
/* A weighted graph has vertices 0, 1, .. to store  
   data of type T and edges which connect vertices.  
   Every edge has a weight of type double. */
```

and there is an additional operation:

```
abstract double get_weight  
  ( graph g, size_t i, size_t j );  
precondition: 0 <= i, j < number_of_vertices(g)  
  and is_edge(g, i, j) is true;  
postcondition: get_weight(g, i, j) is the weight  
  of the edge (i, j);
```

A more general ADT for a weighted graph would leave the type of the weight as defined by a template.

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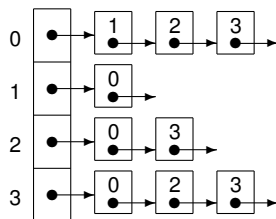
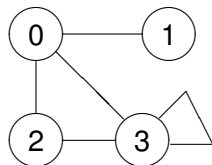
- adjacency matrix
- adjacency list

# adjacency list representation

An adjacency list representation

- uses an array of lists, one list per vertex;
- the list  $i$  stores the vertices adjacent to vertex  $i$ .

An example:



The vertices in the list need not to be ordered.

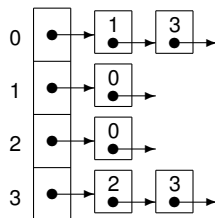
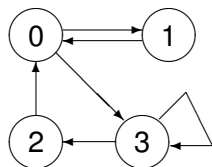
The number of elements in list  $i$  equals the degree of vertex  $i$ .

# adjacency list representation for a directed graph

An adjacency list representation

- uses an array of lists, one list per vertex;
- the list  $i$  stores the vertices  $j$  for all edges  $(i, j)$ .

An example:



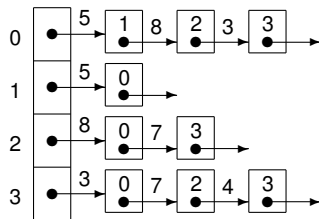
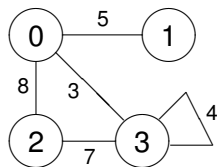
The vertices in the list need not to be ordered.

# adjacency list representation for a weighted graph

An adjacency list representation

- uses an array of lists, one list per vertex;
- the list  $i$  stores the vertices adjacent to vertex  $i$ ;
- the weight of an edge is stored with the pointer to the next vertex.

An example:



The vertices in the list need not to be ordered.

Exercise: draw the adjacency list for a weighted digraph.

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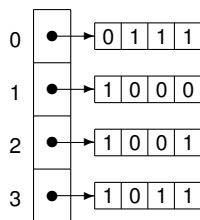
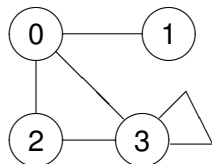
- adjacency matrix
- adjacency list

# adjacency matrix representation

An adjacency matrix representation

- uses an array of arrays, one array per vertex;
- the  $i$ -th array stores a boolean at its  $j$ -th position:
  - ▶ the  $j$ -th boolean in array  $i$  is false if no edge from  $i$  to  $j$ ;
  - ▶ the  $j$ -th boolean in array  $i$  is true if an edge from  $i$  to  $j$ .

An example:



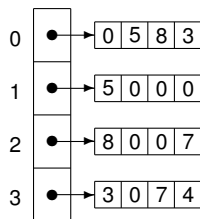
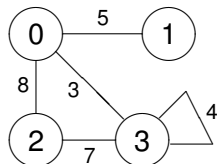
Exercise: draw the adjacency matrix for a digraph.

# adjacency matrix representation for a weighted graph

An adjacency matrix representation

- uses an array of arrays, one array per vertex;
- the  $i$ -th array stores the weight at its  $j$ -th position:
  - ▶ the  $j$ -th number in array  $i$  is zero if no edge from  $i$  to  $j$ ;
  - ▶ the  $j$ -th number in array  $i$  is the weight if an edge from  $i$  to  $j$ .

An example:



Exercise: draw the adjacency matrix for a weighted digraph.

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## graph as adjacency matrix

For the array of arrays, we use the STL `vector` class.

An adjacency matrix is of type `vector< vector<bool> >`.

For a weighted graph, we use `vector< vector<double> >`.

We have four types:

- 1 undirected, unweighted
- 2 directed, unweighted
- 3 undirected, weighted
- 4 directed, weighted

We generate random matrices to test the data structures.

## prototypes for directed, unweighted graphs

```
void allocate_boolean_matrix
    ( vector< vector<bool> >& mat, size_t dim );
// Allocates space for an adjacency matrix
// of dimension dim.

void random_directed_graph
    ( vector< vector<bool> >& mat );
// Fills up the adjacency matrix for a random
// directed graph.

void write_boolean_matrix
    ( vector< vector<bool> >& mat );
// Writes the matrix to screen.
```

## allocation of the matrix

```
void allocate_boolean_matrix
( vector< vector<bool> >& mat, size_t dim )
{
    for(size_t i=0; i<dim; i++)
    {
        vector<bool> row;

        for(size_t j=0; j<dim; j++)
            row.push_back(0);

        mat.push_back(row);
    }
}
```

## generating a random graph

```
void random_directed_graph
( vector< vector<bool> >& mat )
{
    const size_t dim = mat.size();

    for(size_t i=0; i<dim; i++)
        for(size_t j=0; j<dim; j++)
            mat[i][j] = (rand() % 2);
}
```

## writing the matrix

```
void write_boolean_matrix
( vector< vector<bool> >& mat )
{
    const size_t dim = mat.size();

    for(size_t i=0; i<dim; i++)
    {
        for(size_t j=0; j<dim; j++)
            cout << " " << mat[i][j];
        cout << endl;
    }
}
```

## a random symmetric matrix

In a symmetric matrix, the lower half equals the upper half.

```
void random_undirected_graph
( vector< vector<bool> >& mat )
{
    const size_t dim = mat.size();

    for(size_t i=0; i<dim; i++)
    {
        for(size_t j=0; j<i; j++) // copy lower half
            mat[i][j] = mat[j][i];

        for(size_t j=i; j<dim; j++) // generate upper
            mat[i][j] = (rand() % 2);
    }
}
```

## weights in the interval [0.1, 9.9].

```
void random_directed_weighted_graph
( vector< vector<double> >& mat )
{
    const size_t dim = mat.size();
    for(size_t i=0; i<dim; i++)
        for(size_t j=0; j<dim; j++)
        {
            size_t rnd = rand() % 2; // edge or not
            if(rnd == 0)
                mat[i][j] = 0.0; // no edge => 0 weight
            else
            {
                rnd = 1 + rand() % 99; // random weight
                mat[i][j] = ((double) rnd)/10.0;
            }
        }
}
```

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# graph as adjacency list

We use the STL `vector` and `list` classes.

The array of lists is then of type `vector< list<size_t> >`.

For a weighted graph,

we use `vector< list< pair<size_t, double> > >`.

Memory allocation happens in two stages:

- 1 The vector is allocated immediately after its declaration.
- 2 Space in the list is allocated for each new edge.

## prototypes for directed, unweighted graphs

```
void allocate_adjacency_list  
    ( vector< list<size_t> >& lst, size_t dim );  
// Allocates space for an adjacency list  
// of dimension dim.
```

```
void random_directed_graph  
    ( vector< list<size_t> >& lst );  
// Fills up the adjacency list for a random  
// directed graph.
```

```
void write_adjacency_list  
    ( vector< list<size_t> >& lst );  
// Writes the adjacency list to screen.
```

## allocating space

```
void allocate_adjacency_list
( vector< list<size_t> >& lst, size_t dim )
{
    for(size_t i=0; i<dim; i++)
    {
        list<size_t> row;
        lst.push_back(row);
    }
}
```

## generating a directed, unweighted graph

```
void random_directed_graph
( vector< list<size_t> >& lst )
{
    const size_t dim = lst.size();

    for(size_t i=0; i<dim; i++)
    {
        for(size_t j=0; j<dim; j++)
        {
            size_t rnd = rand() % 2;
            if(rnd != 0) lst[i].push_back(j);
        }
    }
}
```

## writing the adjacency list

```
void write_adjacency_list
( vector< list<size_t> >& lst )
{
    const size_t dim = lst.size();

    for(size_t i=0; i<dim; i++)
    {
        cout << i;

        for(list<size_t>::const_iterator it=lst[i].begin();
            it != lst[i].end(); it++)
            cout << " -> " << *it;

        cout << endl;
    }
}
```

## summary + exercises

Graphs representations are adjacency lists or adjacency matrices.

### Exercises:

- 1 For the Graph ADT, write a class definition where the private data attribute is an adjacency matrix.
- 2 Give the implementation of the class of the previous exercise.
- 3 For the Graph ADT, write a class definition where the private data attribute is an adjacency list.
- 4 Give the implementation of the class of the previous exercise.