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#### MCS 572 Lecture 15 Introduction to Supercomputing Jan Verschelde, 30 September 2024

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## mapping jobs to processors

In parallel shared memory computing, we apply the work crew model.

We distinguish between static and dynamic work assignment:

<sup>1</sup> Static: before the execution of the program.

*Each worker has its own queue of jobs to process.*

- $+$  Ideal speedup if jobs are evenly distributed,
- − if one worker gets all long jobs, then unbalanced.
- <sup>2</sup> Dynamic: during the execution of the program.

*Workers process the same queue of jobs.*

- $+$  The size of each job is taken into account,
- − synchronization overhead may dominate for small jobs and when there are many workers.

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### task based programming and work stealing

Tasks are much lighter than threads.

- starting and terminating a task is much times faster than starting and terminating a thread; and
- a thread has its own process id and own resources, whereas a task is typically a small routine.

In scheduling threads on processors,

we distinguish between work sharing and work stealing:

- In work sharing, the scheduler attempts to migrate threads to under-utilized processors in order to distribute the work.
- In work stealing, under-utilized processors attempt to steal threads from other processors.

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# work stealing as hybrid work assignment

Work stealing is illustrated

as a hybrid between static and dynamic work assignment:

- **1** Each worker starts with its own queue.
- 2 An idle worker will work on jobs of other queues.

Main benefit over dynamic work assignment: synchronization overhead occurs only at the end of the execution.

### setup of the Julia program

**1** As many queues as the number of threads are generated:

- $\triangleright$  even indexed queues have small jobs,
- $\triangleright$  odd indexed queues have large jobs.

This generates unbalanced job queues to test the work stealing.

- <sup>2</sup> The *i*-th worker starts processing the *i*-th job queue.
- <sup>3</sup> Every queue has an index to the current job. In Julia, this index is of type  $\text{Atomic} \{ \text{Int} \}$ , for mutual exclusive access.
- <sup>4</sup> After the *i*-th worker is done with its *i*-th job queue, it searches for jobs over all *j*-th queues, for  $j \neq i$ .

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### making the job queues

```
using Base.Threads
```

```
nt = nthreads()
```

```
nbr = 10 # number of jobs in each queue
# allocate memory for all job queues
jobs = [zeros(nbr) for i=1:nt]
```

```
# every worker generates its own job queue
# even indexed queues have light work loads
@threads for i=1:nt
    if i % 2 == 0jobs[i] = rand((1, 2, 3), nbr)else
         i\text{obs}[i] = \text{rand}((4, 5, 6), \text{nbr})end
    println("Worker ", threadid(), " has jobs ",
             i\deltaiobs[i], " ", sum(i\deltaobs[i]))
```
end

# running the program

Each number in the job queue represents the time each job takes.



The ... represents omitted numbers for brevity.

The last number of the output is the sum of the times of the jobs.

Workers 2 and 4 has clearly lighter loads, compared to workers 1 and 3.

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### every worker starts processing its own queue

```
jobidx = [Atomic{Int}(1) for i=1:nt]@threads for i=1:nt
    while true
        myjob = atomic\_add! (jobidx[i], 1)if myjob > length(jobs[i])
            break
        end
        println("Worker ", threadid(),
                 " spends ", jobs[i][myjob], " seconds",
                 " on job ", myjob, " \ldots")
        sleep(jobs[i][myjob])
        jobs[i][myjob] = threadid()
    end
```
Observe the use of the Atomic{Int} for the indices.

The myjob = atomic\_add!(jobidx[i], 1)

- $\bullet$  increments the  $\text{jobidx}[i]$  after returning its value.
- **•** This statement is executed in a critical section.

### idle threads steal work

```
println("Worker ", threadid(), " will steal jobs ...")
more2steal = true
while more2steal
    more2<sub>st</sub>ea1 = falsefor i=1:threadid() -1myjob = atomic\_add! (jobidx[j], 1)if myjob \leq length(jobs[j])
            println("Worker ", threadid(),
                     " spends ", jobs[j][myjob], " seconds",
                     " on job ", myjob, " of ", j, " \dots")
            sleep(jobs[j][myjob])
             iobs[i][myiob] = threadid()
        end
        more2steal = (myjob < length(jobs[j]))end
    for j=threadid()+1:nt # is similar to previous code
```
#### an example of an output

Worker 4 spends 1.0 seconds on job 7 ... Worker 2 will steal jobs ... Worker 2 spends 4.0 seconds on job 4 of 1 ... Worker 4 spends 3.0 seconds on job 8 ... Worker 1 spends 4.0 seconds on job 5 ... Worker 4 spends 3.0 seconds on job 9 ... Worker 2 spends 4.0 seconds on job 5 of 3 ... Worker 3 spends 5.0 seconds on job 6 ... Worker 4 spends 3.0 seconds on job 10 ... Worker 1 spends 6.0 seconds on job 6 ... Worker 3 spends 6.0 seconds on job 7 ... Worker 4 will steal jobs ... Worker 4 spends 6.0 seconds on job 7 of 1 ... Worker 1 spends 4.0 seconds on job 8 ...

Worker 2 is done first, takes job 4 of worker 1. Worker 1 then continues with job 5. When worker 4 is done, it takes job 7 of worker 1. Worker 1 then continues with job 8.

### discussion

- Implementing a work crew with work stealing is not much more complicated than dynamic load balancing.
- The idle workers start at the first queue and then progress linearly, which may be good if the first queue contains all important jobs.
- In an alternative work stealing scheme, idle workers would start in the queue of their immediate neighbors.

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### Numba for multithreading in Python

- Numba is an open-source JIT compiler that translates a subset of Python and NumPy into fast machine code using LLVM, via the llvmlite Python package.
- It offers a range of options for parallelising Python code for CPUs and GPUs, often with only minor code changes.
- Started by Travis Oliphant in 2012, under active development at https://github.com/numba/numba.
- **To use, do** pip install numba.
- The example on the next slide works on Windows.

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### an example from the wikipedia page

```
import numba
import random
@numba.jit
def monte_carlo_pi(n_samples: int) -> float:
    "" "
    Applies Monte Carlo to estimate pi.
    "" "
    acc = 0for i in range(n_samples):
        x = random.random()y = random.random()if (x**2 + y**2) < 1.0:
            acc += 1return 4.0 \times acc / n samples
p = monte_carlo_pi(1000000)
print(p)
```
# <span id="page-17-0"></span>Work Stealing

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# Parsl: Parallel Scripting in Python

- Parsl provides an intuitive, pythonic way of parallelizing codes by annotating "apps": Python functions or external applications that run concurrently.
- **Parsl works seamlessly with Jupyter notebooks.**
- Write once, run anywhere. From laptops to supercomputers.
- **To use, do** pip install parsl.
- **•** The example on the next slide was executed on WSL, Window Subsystem for Linux, Ubuntu 22.04.

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#### an example from the parsl user guide

```
from parsl import python_app
  import parsl
  parsl.load()
  # Map function that returns double the input integer
  @python_app
  def app_double(x):
      return x*2
  # Reduce function that returns the sum of a list
  @python_app
  def app_sum(inputs=()):
      return sum(inputs)
  # Create a list of integers
  items = range(0, 4)# Map phase: apply the double *app* function to each item in list
  mapped_results = []
  for i in items:
      x = app\_double(i)mapped_results.append(x)
  # Reduce phase: apply the sum *app* function to the set of results
  total = appsum(inputs=mapped results)print(total.result())
                                                                (0.125 \times 10^{-14} \text{ m}) \times 10^{-14} \text{ m}QQIntroduction to Supercomputing (MCS 572) Work Stealing L-15 30 September 2024 20 / 46
```
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# the Intel Threading Building Blocks

The Intel TBB is a library that helps you leverage multicore performance *without having to be a threading expert*.

The advantage of Intel TBB is that it works at a higher level than raw threads, yet does not require exotic languages or compilers.

The library differs from others in the following ways:

- TBB enables you to specify logical parallelism instead of threads;
- TBB targets threading for performance;
- **TBB is compatible with other threading packages:**
- TBB emphasizes scalable, data parallel programming;
- $\bullet$  TBB relies on generic programming, (e.g.: use of STL in C $++$ ).

Open Source, download at

http://threadingbuildingblocks.org/ which redirects to a github page.

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### oneTBB and oneAPI

TBB is part of oneAPI which aims to offer one single programming model for CPU, GPU, FPGA accelerators.

On M1 MacBook, the openAPI installs with brew, via brew install tbb and brew install onedpl where onedpl is formerly known as parallelstl a C++ Standard library algorithms with support for execution policies.

To use TBB, look on your system for the location of header files and the libraries.

# saying hello

```
#include <cstdio>
#include <tbb.h>
using namespace tbb;
class say_hello
{
   const char* id:
   public:
      say_hello(const char* s) : id(s) { }
      void operator( ) ( ) const
      {
         printf("hello from task %s\n",id);
      }
};
```
A class in C++ is a like a struct in C for holding data attributes and functions (called methods).

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# the main function

```
int main( )
{
  task_group tg;
  tg.run(say hello("1")); // spawn 1st task and return
  tg.run(say hello("2")); // spawn 2nd task and return
  tg.wait( ); // wait for tasks to complete
}
```
The run method spawns the task immediately, but does not block the calling task, so control returns immediately.

To wait for the child tasks to finish, the classing task calls wait.

Observe the syntactic simplicity of task group.

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### raising complex numbers to a large power

Consider the following problem:

Input: 
$$
n \in \mathbb{Z}_{>0}
$$
,  $d \in \mathbb{Z}_{>0}$ ,  $\mathbf{x} \in \mathbb{C}^n$ .  
Output:  $\mathbf{y} \in \mathbb{C}^n$ ,  $y_k = x_k^d$ , for  $k = 1, 2, ..., n$ .

To avoid overflow, we take complex numbers on the unit circle.

In C++, complex numbers are defined as a template class. To instantiate the class complex with the type double we declare

#include <complex>

using namespace std;

typedef complex<double> dcmplx;

#### random complex doubles

```
#include <cstdlib>
#include <cmath>
```
dcmplx random dcmplx ( void ); // generates a random complex number // on the complex unit circle

We compute  $e^{2\pi i \theta} = \cos(2\pi \theta) + i \sin(2\pi \theta),$  for random  $\theta \in [0,1]$ :

```
dcmplx random dcmplx ( void )
{
   int r = \text{rand}();
   double d = ((double) r)/RAND MAX;double e = 2*M PI*d;
   dcmplx c(cos(e),sin(e));
   return c;
}
```
### writing arrays

```
#include <iostream>
#include <iomanip>
```

```
void write numbers ( int n, dcmplx *x );
// writes the array of n doubles in x
```

```
Observe the local declaration int in the for loop,
the scientific formatting, and the methods real() and imag():
```

```
void write numbers ( int n, dcmplx *x )
{
   for(int i=0; i<n; i++)
      cout << scientific << setprecision(4)
           << "x[" << i << "] = ( " << x[i].real()
           << ", " << x[i].imag() << ") \n";
```
}

### computing powers

```
void compute powers ( int n, dcmplx *x,
                      dcmplx *y, int d );
// for arrays x and y of length n,
// on return y[i] equals x[i]**d
```
The plain  $for(int \text{ int})$  loop avoids repeated squaring:

```
void compute powers ( int n, dcmplx *x,
                    dcmplx *y, int d)
{
  for(int i=0; i < n; i++) // y[i] = pow(x[i], d);{ // pow is too efficient
     dcmplx r(1.0, 0.0);
     for(int j=0; j < d; j++) r = r*x[i];
     y[i] = r;}
}
```
#### command line arguments

```
$ ./powers serial
how many numbers ? 2
x[0] = (-7.4316e-02, 9.9723e-01)x[1] = (-9.0230e-01, 4.3111e-01)give the power : 3
x[0] = (2.2131e-01, -9.7520e-01)x[1] = (-2.3152e-01, 9.7283e-01)
```
\$ ./powers\_serial 2 3 1 x[0] = ( -7.4316e-02 , 9.9723e-01) x[1] = ( -9.0230e-01 , 4.3111e-01) x[0] = ( 2.2131e-01 , -9.7520e-01) x[1] = ( -2.3152e-01 , 9.7283e-01)

\$ time ./powers\_serial 1000 1000000 0

real 0m20.139s

user 0m20.101s

sys 0m0.000s

# the main program

```
int main ( int argc, char *argv[] )
{
   int v = 1; // verbose if > 0
   if(argc > 3) v = atoi(aray[3]);
   int dim; // get the dimension
   if(argc > 1)
      dim = atoikaray[1]);
   else
   {
      cout << "how many numbers ? ";
      cin >> dim;
   }
   // fix the seed for comparisons
   srand(20120203); //srand(time(0));
   dcmplx r[dim];
   for(int i=0; i<dim; i++)
      r[i] = random_dcmplx();
   if(v > 0) write numbers(dim,r);
                                               A \equiv A \equiv A B \equiv A
```
# the main program continued

```
int deg; // get the degree
if(argc > 1)
  deq = atoi(arqv[2]);
else
{
  cout << "give the power : ";
  cin >> deg;
}
dcmplx s[dim];
compute powers(dim,r,s,deg);
if(v > 0) write numbers(dim,s);
return 0;
```
}

D.

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### the speedup

\$ time ./powers\_serial 1000 1000000 0

- real 0m20.139s
- user 0m20.101s
- sys 0m0.000s
- \$ time ./powers tbb 1000 1000000 0
- real 0m1.191s
- user 0m35.170s
- sys 0m0.043s

# The speedup:  $\frac{20.139}{1.191} = 16.909$  on two 8-core CPUs.

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### on M1 MacBook Air

- % /usr/bin/time /tmp/powers\_serial 1000 1000000 0 16.23 real 13.15 user 0.02 sys
- % /usr/bin/time /tmp/powers\_tbb 1000 1000000 0 2.30 real 16.96 user 0.04 sys

### the class ComputePowers

```
class ComputePowers
{
   dcmplx *const c; \frac{7}{1} numbers on input<br>int d: \frac{7}{1} degree
                   // degree
   dcmplx *result; // output
   public:
      ComputePowers(dcmplx x[], int deg, dcmplx y[])
          : c(x), d(deq), result(y) { }
      void operator()
          ( const blocked_range<size_t>& r ) const
       {
          for(size t i=r.begin(); i!=r.end(); ++i)
          {
             dcmplx z(1.0, 0.0);
              for(int j=0; j < d; j++) z = z * c[i];
             result[i] = z;}
       }
};
                                                   A EXA EXAMPLE VOOR
```
#### tbb/blocked\_range.h

#include "tbb/blocked\_range.h"

template<typename Value> class blocked\_range

A blocked\_range represents a half open range [*i*, *j*) that can be recursively split.

```
void operator()
   ( const blocked_range<size_t>& r ) const
{
   for(size t i=r.begin(); i!=r.end(); ++i)
   {
```
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### calling the parallel\_for

#include "tbb/tbb.h" #include "tbb/blocked\_range.h" #include "tbb/parallel for.h"

using namespace tbb;

One line changes in the main program:

```
parallel_for(blocked_range<size_t>(0,dim),
             ComputePowers(r,deg,s));
```
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# an application of work stealing



#### from the Intel Threading Building Blocks Tutorial



# what if no worker is available?



No Available Worker

#### from the Intel Threading Building Blocks Tutorial

### the class SumIntegers

```
class SumIntegers
{
   int *data;
   public:
      int sum;
      SumIntegers ( int *d ) : data(d), sum(0) {}
      void operator()
         ( const blocked_range<size_t>& r )
      {
         int s = sum; // must accumulate !
         int *d = data;size t end = r.end();
         for(size t i=r.begin(); i != end; ++i)
            s \neq d[i];sum = s;
      }
```
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### split and join methods

```
// the splitting constructor
      SumIntegers ( SumIntegers& x, split ) :
         data(x.data), sum(0) \{\}// the join method does the merge
      void join ( const SumIntegers& x ) { sum += x.sum; }
};
int ParallelSum ( int *x, size_t n )
   SumIntegers S(x);
  parallel reduce(blocked range<size t>(0,n), S);
   return S.sum;
}
```
{

### code in the main program

```
int *d;
d = (int*)calloc(n, sizeof(int));
for(int i=0; i<n; i++) d[i] = i+1;
int s = ParallelSum(d, n);
```
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### <span id="page-45-0"></span>Exercises

- <sup>1</sup> A permanent is similar to a determinant but then without the alternating signs. Develop a task-based parallel program to compute the permanent of a 0/1 matrix. Why is work stealing appropriate for this problem?
- <sup>2</sup> Modify the hello world! program with TBB so that the user is first prompted for a name. Two tasks are spawned and they use the given name in their greeting.
- Modify powers tbb.cpp so that the *i*th entry is raised to the power *d* − *i*. In this way not all entries require the same work load. Run the modified program and compare the speedup to check the performance of the automatic task scheduler.