# Some solutions for the math 125 final review packet 

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The students in my section decided via democratic process the problems in the Math 125 final exam practice packet for which they wanted to see solutions. I this guide I do not write out the statements of the problems, so one will need a copy of the Math 125 final review (Click in the final review button at the Math 125 Exams page.) I have no doubt that this document contains typos or mistakes; please let me of these. Good luck, and study well.
4. Mine 1 produces 4 tons of iron and 1 ton of copper at $\$ 3,200$ a day, and mine 2 produces 2 tons of iron and 2 tons of copper at $\$ 2,600$ a day. We need at least 40 tons of iron and 18 tons of copper. Let x be the number of days running mine 1 and $y$ be the number of days running mine 2 . Then we have the following system.

$$
\left\{\begin{array}{l}
\text { Minimize } 3,200 x+2,600 y \\
4 x+2 y \geq 40 \\
2 x+2 y \geq 18 \\
x, y \geq 0
\end{array}\right.
$$

6.     - 

$$
\left(\begin{array}{ccccc}
3 & 2 & -2 & -6 & 5 \\
-3 & 12 & 7 & -2 & -6 \\
0 & -4 & 5 & 1 & 3 \\
1 & -2 & -3 & -7 & 5
\end{array}\right)\left(R_{4} \rightarrow 4 R_{4}\right) \rightarrow\left(\begin{array}{ccccc}
3 & 2 & -2 & -6 & 5 \\
-3 & 12 & 7 & -2 & -6 \\
0 & -4 & 5 & 1 & 3 \\
4 & -8 & -12 & -28 & 20
\end{array}\right)
$$

- 

$$
\left(\begin{array}{ccccc}
3 & 2 & -2 & -6 & 5 \\
-3 & 12 & 7 & -2 & -6 \\
0 & -4 & 5 & 1 & 3 \\
1 & -2 & -3 & -7 & 5
\end{array}\right)\left(R_{1} \rightarrow-2 R_{4}+R_{1}\right) \rightarrow\left(\begin{array}{ccccc}
1 & 6 & 4 & 8 & -5 \\
-3 & 12 & 7 & -2 & -6 \\
0 & -4 & 5 & 1 & 3 \\
1 & -2 & -3 & -7 & 5
\end{array}\right)
$$

- 

$$
\left(\begin{array}{ccccc}
3 & 2 & -2 & -6 & 5 \\
-3 & 12 & 7 & -2 & -6 \\
0 & -4 & 5 & 1 & 3 \\
1 & -2 & -3 & -7 & 5
\end{array}\right)\left(R_{3} \rightarrow 5 R_{1}+R_{3}\right) \rightarrow\left(\begin{array}{ccccc}
3 & 2 & -2 & -6 & 5 \\
-3 & 12 & 7 & -2 & -6 \\
15 & 6 & -5 & -29 & 28 \\
1 & -2 & -3 & -7 & 5
\end{array}\right)
$$

11. 

|  | Metal | Plastic | Wood |
| :---: | :---: | :---: | :---: |
| Metal | 0.08 | 0 | 0.05 |
| Plastic | 0.03 | 0.11 | 0.04 |
| Wood | 0.07 | 0.02 | 0.10 |

Thus, we have that

$$
A=\left(\begin{array}{ccc}
0.08 & 0 & 0.05 \\
0.03 & 0.11 & 0.04 \\
0.07 & 0.02 & 0.10
\end{array}\right)
$$

We are given that $D=\left(\begin{array}{l}32 \\ 25 \\ 40\end{array}\right)$ where the numbers in $D$ denote millions of dollars. Now on your calculator, compute $(I-A)^{-1} D$. Round this to the hundredths place. Remember that you must multiply with $D$ on the right.
13. Let $x$ be the number of package $A$, and $y$ be the number of package B. From the statement of the problem we see that or linear programming question is the following.

$$
\left\{\begin{array}{l}
\text { Maximize } 175 x+240 y \\
140 x+185 y \leq 6,000 \\
3 x+5 y \leq 300 \\
x, y \geq 0
\end{array}\right.
$$

14. Let x be the number of tablets and y be the total cost. This question gives us two linear equations. For the Thailand plant we have $y=79 x+9,000$ and for the Malaysia plant we have $y=71 x+9,400$.

- This question is just asking us to set $x=600$ and see which function has a lesser value. In particular the costs at Thailand are $\$ 79 \cdot 600+\$ 9,000=\$ 56,400$ and similarly the costs at Malaysia are $\$ 52,000$. Thus, it's cheaper to produce in Malaysia.
- Let $x_{1}$ be the number of units made in Thailand and $x_{2}$ be the number of units made in Malaysia. Since we have to produce 1000 units we see that $x_{1}=1000-x_{2}$. All that is left to do is set the two linear equations to be equal and plug in $1000-x_{2}$ for $x_{1}$ in the Thailand equation. Solving $79\left(1000-x_{2}\right)+9,000=71 x_{2}+9,400$ for $x_{2}$ we have that $x_{2}=524(524$ tablets must be made in Malaysia), and thus $x_{1}=476$ ( 476 tablets must be made in Thailand).

15. We use the data given to set up the following table.

|  | Millions of people | Number of stoves sold |
| :---: | :---: | :---: |
| Atlanta | 6.1 | 13,286 |
| Tampa | 2.8 | 5,123 |
| Miami | 6.4 | 17,522 |
| Charlotte | 2.5 | 4,848 |
| Greenville | 1.4 | 3,613 |

Using the LinReg feature we get that the line of best fit is $y=2647 x-1286$. ${ }^{1}$

- Just plug in $x=2.9$ to the equation we just found.
- Solving $20,000=2647 x-1286$ for $x$ gives us that we need about $x=8.04$ million people.

[^0]17. We need a linear programming problem (LPP) in two variables, but at first lets start with a LPP in three. Let $L$ be the number of lions, $B$ be the number of bears, and $T$ be the number of tigers. Reading the question carefully will give us the following LPP in three variables.
\[

\left\{$$
\begin{array}{l}
\text { Minimize 40L+30T+25B } \\
L+B+T=150 \\
2 L \geq T \\
B \leq T \\
B, L, T \geq 25
\end{array}
$$\right.
\]

However, we notice that we can solve the equation $L+B+T=150$ for one of the variables. You may pick you favorite one, but I'm solving the question right now and my favorite one is $L$. So we have $L=150-B-T$, and we can substitute this $L$ in the LPP to reduce the number of variables.

$$
\left\{\begin{array}{l}
\text { Minimize } 40(150-\mathrm{B}-\mathrm{T})+30 \mathrm{~T}+25 \mathrm{~B} \\
(150-B-T)+B+T=150 \\
2(150-B-T) \geq T \\
B \leq T \\
B,(150-B-T), T \geq 25
\end{array}\right.
$$

This is a mess. Note the the second equality reduces to $150=150$, so we discard that equality now. Also the last row gives is that $(150-B-T) \geq 25 \Rightarrow B+T \leq 125$ which is a non-trival inequality. We need to include in in our new LPP. Now just clean up the rest of the inequalities to get

$$
\left\{\begin{array}{l}
\text { Minimize } 6000-15 \mathrm{~B}-10 \mathrm{~T} \\
B+T \leq 125 \\
2 B+3 T \leq 300 \\
B-T \leq 0 \\
B, T \geq 25
\end{array}\right.
$$

19.     - The first constraint is changing from 34 to 30 , so in this case $h=-4=30-34$. As written, the first constraint corresponds to the $u$ column, so the new right hand column is

$$
\left(\begin{array}{c}
2-4(-5 / 2) \\
4-4(1) \\
13-4(-1 / 2) \\
47-4(1 / 2)
\end{array}\right)=\left(\begin{array}{c}
12 \\
0 \\
15 \\
45
\end{array}\right)
$$

- Mostly the same story. 30 is changing to 32 so $h=2=32-30$. Also, the third constraint corresponds to the w column so the new right hand side is

$$
\left(\begin{array}{c}
2+2(1) \\
4+2(-1) \\
13+2(1) \\
47+2(1)
\end{array}\right)=\left(\begin{array}{c}
4 \\
2 \\
15 \\
49
\end{array}\right)
$$

- In this case, we have that 34 is changing to $34+\mathrm{h}$, as above we just look at the difference of $34+h-34=$ $h$. Constraint 1 corresponds to the $u$ column so we have

$$
\left(\begin{array}{c}
2+h(-5 / 2) \\
4+h(1) \\
13+h(-1 / 2) \\
47+h(1 / 2)
\end{array}\right)
$$

since the $M$ value doesn't matter when we are looking to remain feasible, we only need to look at the following inequalities.

$$
\left\{\begin{array}{l}
2+h(-5 / 2) \geq 0 \\
4+h(1) \geq 0 \\
13+h(-1 / 2) \geq 0
\end{array}\right.
$$

Clean this up to get

$$
\left\{\begin{array}{l}
h \leq 4 / 5 \\
h \geq-4 \\
h \leq 26
\end{array}\right.
$$

To make this even clearer, rewrite this as

$$
\left\{\begin{array}{l}
h \leq 4 / 5 \\
-4 \leq h \\
h \leq 26
\end{array}\right.
$$

We can now see that $-4 \leq h \leq 4 / 5$ if we want to stay feasible.
25. We are given that $A B=C$ and asked to find $b_{2,1}+b_{2,2}$. As we discussed in the discussion sections this just means we have to find the entries in the second row first column and the second row second column in B and add them together. Since we know $A$ and $C$ and we see that the determinant of $A$ is not 0 , we know that we can invert $A$. This gives us that

$$
B=A^{-1} C
$$

Using either your calculator or doing the calculation by hand, we will see that $A^{-1}=\left(\begin{array}{cc}2 & 1 \\ 5 / 2 & 3 / 2\end{array}\right)$ Thus

$$
B=A^{-1} C=\left(\begin{array}{cc}
2 & 1 \\
5 / 2 & 3 / 2
\end{array}\right)\left(\begin{array}{cc}
-6 & -2 \\
11 & 6
\end{array}\right)=\left(\begin{array}{cc}
-1 & 2 \\
3 / 2 & 4
\end{array}\right)
$$

The entry in the second row first column is $3 / 2$ and the entry in the second row second column is 4 . So $b_{2,1}+b_{2,2}=3 / 2+4=11 / 2$
29. Let's define the payoff matrix as the money that Player C must pay Player R. The matrix then looks like this

|  |  | Player R |  |  |
| :---: | :---: | :---: | :---: | :---: |
| Player C | $\$ 5$ | $\$ 1$ | $\$ 10$ | $\$ 50$ |
|  | $\$ 20$ | 19 | -5 | -45 |
|  |  |  | 10 | -40 |

Asking if this game is strictly determined is asking if this matrix has saddle point. So we need to perform the following.

1) Determine the least element in each row. (-40,-45)
2) Choose the largest of these elements. (-40)
3) Determine the greatest element in each column. (19, 10, -40)
4) Choose the smallest of these elements. (-40)

If the element chosen in step 2 and 4 is the same (it is) we have found our saddle point.
Following this guide we see that -40 is the saddle point. Therefore, thus game is strictly determined, and the Value of the game is $v=-40$
31. (With David N. Reynolds) We start with the matrix $\left(\begin{array}{cc}0 & -3 \\ -1 & 2\end{array}\right)$. We need to add 4 to make every entry positive. So we have the matrix $\left(\begin{array}{ll}4 & 1 \\ 3 & 6\end{array}\right)$ We now construct the LPP for the optimal strategy for player C.

$$
\begin{aligned}
& \text { Maximize } z_{1}+z_{2} \\
& \left\{\begin{array}{l}
4 z_{1}+z_{2} \leq 1 \\
3 z_{1}+6 z_{2} \leq 1 \\
z_{1}, z_{2} \geq 0
\end{array}\right.
\end{aligned}
$$

Which gives us the following initial tableau.

$$
\left[\begin{array}{ccccc|c}
z_{1} & z_{2} & u_{1} & u_{2} & M & \\
4 & 1 & 1 & 0 & 0 & 1 \\
3 & 6 & 0 & 1.4 & 0 & 1 \\
\hline-1 & -1 & 0 & 0 & 1 & 0
\end{array}\right]
$$

Which will give us the following final tableau.

$$
\left[\begin{array}{ccccc|c}
z_{1} & z_{2} & u_{1} & u_{2} & M & \\
1 & 0 & 2 / 7 & -1 / 21 & 0 & 5 / 21 \\
0 & 1 & -1 / 7 & 4 / 21 & 0 & 1 / 21 \\
0 & 0 & 1 / 7 & 1 / 7 & 1 & 2 / 7
\end{array}\right]
$$

Note that this gives us that $v=\frac{1}{z_{1}+z_{2}}=\frac{1}{5 / 21+1 / 21}=\frac{21}{6}=\frac{7}{2}$. Now do the following computations.

- $z_{1} \cdot v=\frac{5}{21} \cdot \frac{7}{2}=\frac{5}{6}$
- $z_{2} \cdot v=\frac{1}{21} \cdot \frac{7}{2}=\frac{1}{6}$

These first two computations give us the entries for the optimal mixed strategy for C, while these next two give us the entries for the optimal mixed strategy for R.

- $u_{1} \cdot v=\frac{1}{7} \cdot \frac{7}{2}=\frac{1}{2}$
- $u_{2} \cdot v=\frac{1}{7} \cdot \frac{7}{2}=\frac{1}{2}$

So the optimal mixed strategy for $C$ is $\left[\begin{array}{c}\frac{5}{6} \\ \frac{1}{6}\end{array}\right]$ and the optimal mixed strategy for $R$ is $\left[\begin{array}{ll}\frac{1}{2} & \frac{1}{2}\end{array}\right]$ The last step is to remember that the $v$ we found it the value for the matrix we shifted by 4 . Thus, the value of the original matrix is $v-4=\frac{7}{2}-4=-\frac{1}{2}$.


[^0]:    ${ }^{1}$ (The problem says to round to the nearest thousandth, but I don't have a TI calculator right now, this is the idea, and this will be very close to the answer you will get with your calculator.)

