

EXERCISES FOR METAMATH II: THIRD SET

1. EXERCISE

Let $\Sigma_k = \{0, \dots, k-1\}$. Suppose $L \subseteq \Sigma_k^*$ is a regular language and define

$$R = \{x \in L \mid x \text{ is the lexicographically largest element of } L \cap \Sigma_k^{|x|}\}.$$

Show that R is regular.

2. EXERCISE

Let L be the set of all non-square natural numbers written in base 10, i.e.,

$$L = \{2, 3, 5, 6, 7, 8, 10, 11, 12, 13, 14, 15, 17, 18, \dots\}.$$

Find the (finite) set of minimal strings in L with respect to the subsequence ordering \preceq . I.e., find the smallest finite subset $S \subseteq L$ such that for any $x \in L$ there is some $y \in S$ with $y \preceq x$.

3. EXERCISE

Show that the following are both presentations of the trivial group.

- $\langle x, y \mid xyx^{-1} = y^2, yxy^{-1} = x^2 \rangle$,
- $\langle x, y, z \mid xyx^{-1} = y^2, yzy^{-1} = z^2, zxz^{-1} = x^2 \rangle$.

4. EXERCISE

Suppose Σ is a finite alphabet and $L \subseteq \Sigma^*$ is a regular language. For $x \in \Sigma^*$, define

$$\text{tree}_L(x) = \{y \in \Sigma^* \mid \text{cone}_L(xy) \neq \emptyset\}.$$

So $\text{tree}_L(x)$ is a (possibly empty) tree on Σ^* , i.e., a set of strings closed under taking prefixes. Let $[\text{tree}_L(x)]$ denote the set of all *infinite branches* of $\text{tree}_L(x)$, i.e., the set of $\alpha \in \Sigma^\omega$ all of whose finite prefixes belong to $\text{tree}_L(x)$. We say that $\text{tree}_L(x)$ is *well-founded* in case $[\text{tree}_L(x)] = \emptyset$.

- Show that there is a constant K only depending on L such that for any $x \in \Sigma^*$, if $\text{tree}_L(x)$ is well-founded, then

$$|\text{tree}_L(x)| \leq K.$$

5. EXERCISE

Find a characterisation of all regular languages $L \subseteq \Sigma^*$ with the following property: For some K dependent on L and for all n ,

$$|L \cap \Sigma^n| \leq K.$$

6. EXERCISE

Find the \sim_L equivalence classes of the following languages.

- $L = \{w \in \{0, 1\}^* \mid |w|_0 = |w|_1\}$,
- $L = \{\alpha^n b^n c^n \mid n \geq 1\}$.

7. EXERCISE

For every $k \geq 1$, show how to construct a regular language L_k over a finite alphabet Σ_k such that for all n ,

$$|L_k \cap (\Sigma_k)^n| = n^k.$$

8. EXERCISE

Let $\Sigma = \{1, 2, \dots, n\}$. A string $w = a_1 a_2 \dots a_k$ is said to be *bipartite* if there is a partition $A \sqcup B = \{1, \dots, k\}$ such that

$$\sum_{i \in A} a_i = \sum_{i \in B} a_i.$$

(a) Show that there is a constant K depending only on n such that any bipartite string w has a bipartite substring $v \preceq w$ with $|v| \leq K$.

(b) Show that the language of bipartite strings is regular.