Lecture 2

1 Point-line incidences

Let P be a set of m (distinct) points and L be a set of n (distinct) lines in the plane. An incidence is a pair $(p, \ell) \in P \times L$ such that $p \in \ell$. Denote $I(P, L) = \{(p, \ell) \in P \times L : p \in \ell\}$ and denote I(m, n) to be the maximum |I(P, L)| over all choices of an m-element point set P and an n-element set of lines L.

1.1 Upper bound proof of Székely, 1997

A drawing of a graph G is a mapping that assigns vertices to points, and edges to continuous arcs connecting the corresponding vertices. The interior of each arc (edge) does not contain any vertices. A pair of arcs (edges) cross if their interiors have a point in common. The crossing number of a graph G, denoted by cr(G), is the minimum number of pairs of arcs (edges) that cross over all drawings of G.

A graph G is *planar* if and only if it has a drawing with no crossings.

Theorem 1.1 (Euler, 1750's). Let G = (V, E) be a connected planar graph drawn in the plan with |F| faces. Then

$$|V| - |E| + |F| = 2.$$

Note that G may contain multiple edges and loops.

Proof. Induction on |V|. Base case |V| = 1 is trivial. For |V| > 2, let e = uv be an edge connecting two distinct vertices. By contracting e, we reduce the number of vertices and edges by one and the proof is complete by the induction hypothesis.

Corollary 1.2. If G is a simple (no loops or multiple edges) planar graph with n vertices, then $|E(G)| \leq 3n - 6$.

Proof. Let F be the set of faces in a planar drawing of G. Then

$$3|F| \le \sum_{f \in F} |f| = 2|E|$$

which implies $|F| \leq 2|E|/3$. Plugging this into Euler's formula gives

$$2 \le |V| - |E| + (2|E|/3),$$

which implies $|E| \leq 3|V| - 6$.

Corollary 1.3. For any simple graph G = (V, E), $cr(G) \ge |E| - 3|V|$.

Proof. If |E| > 3|V|, and some drawing of G had fewer than |E| - 3|V| crossings, then we can delete one edge from each crossing and obtain a planar graph with more than 3|V| edges. This contradiction completes the proof.

Theorem 1.4 (Crossing Lemma, Ajtai-Chvá tal-Newborn-Szemerédi 1982, Leighton 1983). Let G be a simple graph. Then

$$cr(G) \ge \frac{1}{64} \frac{|E|^3}{|V|^2} - |V|.$$

Proof. (Folklore) Consider any drawing of G = (V, E) in the plane, with n vertices, m edges, and x crossings (pairs of edges that cross). We can assume that $m \ge 4n$, otherwise the statement is trivial. Choose a random subset $V' \subset V$ by picking each vertex $v \in V$ independently with probability p. Let G' = (V', E') be the subgraph induced on V' with n' vertices, m' edges, and x' crossings in the *inherited* (same) drawing. Then $\mathbb{E}[n'] = np$, $\mathbb{E}[m'] = mp^2$, and $E[x'] = xp^4$. By Lemma 1.3 and linearity of expectation, we have

$$\mathbb{E}[x'] \ge \mathbb{E}[m'] - 3\mathbb{E}[n'],$$

which implies

 $xp^4 \ge mp^2 - 3np.$

By setting p = 4n/m (recall $m \ge 4n$), we have

$$x \ge \frac{1}{64} \frac{m^3}{n^2}.$$

Proof of the Szemerédi-Trotter theorem. Consider a set P of m points and a set L of n lines in the plane maximizing I(m, n). We define a drawing of a graph G where the vertices of G are the points from P, and two points p, q (vertices) are connected iff p and q lie on a common line $\ell \in L$ *next* to each other. Since a line with k points defines k-1 edges, our graph G contains I(m, n) - nedges.

We have $cr(G) \leq {n \choose 2} \leq n^2$, since two lines cross at most once. On the other hand, by the Crossing Lemma, we have

$$cr(G) \ge \frac{1}{64} \frac{|E|^3}{m^2} - m = \frac{1}{64} \frac{(I(m,n)-n)^3}{m^2} - m.$$

which implies $I(m, n) \le O(m^{2/3}n^{2/3} + m + n)$.

Theorem 1.5 (Székely 1997). Let P be a set of m points and L be a set of n pseudo lines in the plane, that is, a family of curves in which every two members in L have at most one point in common. Then $I(P,L) \leq O(m^{2/3}n^{2/3} + m + n)$.

References

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