

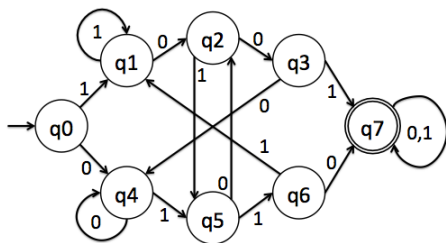
MCS 541 – Computational Complexity  
 Spring 2023  
 Problem Set 1

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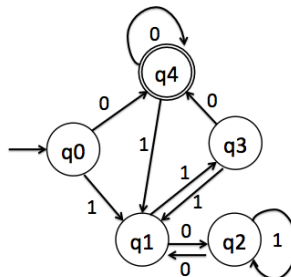
**Due:** 1/23/23 at the beginning of class

1. For the two following DFAs, explain in words the languages  $L_{1,a}$  and  $L_{1,b}$  that they recognize.

a.  $M_{1,a}$



b.  $M_{1,b}$



2. Let  $L_{2,a}$  and  $L_{2,b}$  be regular languages. Show that  $\{x \mid x \in L_{2,a}, x \notin L_{2,b}\}$  is also regular.

3. Let  $L_3 = \{0^n 1^n \mid n \geq 1\}$ . Draw the state diagram for a Turing Machine that decides  $L_3$ . All transitions not depicted will be assumed to go to the reject state, which may or may not be drawn. Note: knowing how to program TMs is not an important, but doing it at least once is a worthwhile exercise.

4. The goal of this problem is to find an uncomputable (or “incomputable”) function. This function,  $S(n)$ , will grow so fast that even Turing machines cannot keep up with it.

a. Let  $L_{4,a} = \{\langle M \rangle \mid M \text{ is a TM and } M \text{ halts on the input } \varepsilon \text{ (i.e. no input)}\}$ . You can assume the input alphabet to be  $\Sigma = \{0, 1\}$ , which is sufficient to encode  $M$ . Is  $L_{4,a}$  recognizable? Is  $L_{4,a}$  decidable? Prove your answers correct.

b. Let  $\mathcal{H}_n$  be the set of  $n$ -state TMs that eventually halt when run on the empty input. Let  $S(n)$  be the maximum number steps a TM in  $\mathcal{H}_n$  can take. Let

$$L_{4,b} = \{m \mid m \in \{0, 1\}^* \text{ s.t. } m = S(n) \text{ for some } n \geq 1\}.$$

Prove that  $L_{4,b}$  is not decidable.

5. Consider the language  $L_5$ , defined as follows:  $L_5 = \begin{cases} \{1\} & \text{if } L_{4,b} \text{ contains infinitely many primes} \\ \{0\} & \text{otherwise.} \end{cases}$   
 Is  $L_5$  decidable? Why or why not?

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<sup>1</sup>Interpret  $m$  as a binary number.