MCS 441 – Theory of Computation I Spring 2013 Problem Set 8

Lev Reyzin

Due: 4/5/13 at the beginning of class

Related reading: Chapters 5 and 6.

Instructions: Atop your problem set, write your name, list your collaborators (see syllabus for the collaboration policy), and indicate whether you are an undergraduate or graduate student.

Rice's Theorem

1. [9 pts] Consider the following theorem.

Theorem 1 (Rice '53) Let P be a language consisting of Turing machine descriptions where P fulfills the following two conditions:

- 1. P contains some but not all TM descriptions.
- 2. For any two TMs M_1 and M_2 , whenever $L(M_1) = L(M_2)$, we have $\langle M_1 \rangle \in P$ iff $\langle M_2 \rangle \in P$.

Then P is undecidable.

Solve the following two problems.

- a. [3 pts] Invent an interesting undecidable language and show it is undecidable using Theorem 1.
- b. [6 pts] Show that both conditions 1. and 2. of Theorem 1 are necessary for it to be true.

Post Correspondence Problem

2. [6 pts] Consider the Post Correspondence Problem (PCP). In Sipser's example on p. 227 ch. 5.2, the strings written on the dominos are over the alphabet $\{a, b, c\}$. That variant of PCP is undecidable. Now consider a different variant of the problem where the alphabet is simply $\{a\}$, i.e. the strings on the dominos are composed of a's. Prove this latter variant of PCP is decidable.

Many-to-one and Turing reductions

- 3. [8 pts] Let $\Sigma = \{0, 1\}$ and $L \subseteq \Sigma^*$.
 - a. [4 pts] Show that L is decidable if and only if $L \leq_T \Sigma^*$.
 - b. [4 pts] Is the statement above true if we replace \leq_T with \leq_m ? Why or why not?

¹To solve this problem, it is not necessary to understand the proof of Theorem 5.15

Kolmogorov complexity

- **4.** [11 pts] Let K(x) be the Kolmogorov complexity (or "descriptive complexity") of string x.
 - a. [6 pts] Show that an oracle Turing machine with an oracle for

$$HALT_{TM} = \{ \langle M, w \rangle \mid M \text{ is a Turing machine that halts on } w \}$$

can compute K(x). Formalize this problem as needed.

b. [5 pts] For $i \ge 1$, give the best upper bound you can for $K(x^i)$ as a function of K(x) and i.